

# A BIJECTION BETWEEN TERNARY TREES AND A SUBCLASS OF MOTZKIN PATHS

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ABSTRACT. A bijection between ternary trees with  $n$  nodes and a subclass of Motzkin paths of length  $3n$  is given. This bijection can then be generalized to  $t$ -ary trees.

## 1. INTRODUCTION

A recent question in the International Mathematics Competition proposed by Petrov and Vershik [4] counts the number of allowed paths from  $(0, 0, 0)$  to  $(n, n, n)$  of a frog that makes steps of length one along the lattice

$$\Omega = \{(x, y, z) \in \mathbb{Z}^3 \mid 0 \leq z \leq y \leq x \leq y + 1\}$$

in exactly  $3n$  moves.

Clearly there are  $n$  steps in each of the three possible directions, and we model each step as follows:

$(x, y, z)$	$(0, 0, 1)$	$(0, 1, 0)$	$(1, 0, 0)$
Step	$\setminus$	$/$	$-$

This along with the restriction  $0 \leq z \leq y \leq x \leq y + 1$ , gives rise to the subclass of Motzkin paths defined below.

**Definition 1.** *An **S-Motzkin path** is a Motzkin path with  $n$  of each type of step such that the following conditions hold*

- *The initial step must be  $-$ ,*
- *between every two  $-$  there is exactly one  $/$ ,*
- *the  $k$ -th occurring  $\setminus$  must occur after at least  $k$  pairs of  $-$  and  $/$ .*

The total number of such paths is  $\frac{1}{2n+1} \binom{3n}{n}$  which is equal to the number of ternary trees with  $n$  nodes [1]. We first provide a mapping from S-Motzkin paths to ternary trees, and then provide the inverse mapping, thus showing that S-Motzkin paths are bijective to ternary trees as well as other combinatorial objects found in [2, 3, 5]. For completeness, an instructive example is given along with a table for  $n = 3$ .

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## 2. BIJECTION

2.1. **S-Motzkin paths to ternary trees.** We define  $\emptyset$  to be the empty path. For an arbitrary S-Motzkin path  $\mathcal{M}$ , the canonical decomposition is

$$\Phi(\mathcal{M}) = (\mathcal{A}, \mathcal{B}, \mathcal{C}),$$

where  $\mathcal{A}$ ,  $\mathcal{B}$ , and  $\mathcal{C}$  represent paths at the left, middle, and right subtrees respectively. Furthermore,

- $\mathcal{C}$  is the path from the penultimate return of the path to the last return, with the initial and last steps removed,
- $\mathcal{A}$  is the path from  $y$  to  $x$  (not including  $x$ ), where  $x$  is the first  $\_$  to the left of  $\mathcal{C}$ , and  $y$  the farthest away  $\_$  from  $x$  such that the path from  $y$  to  $x$  is still a Motzkin path, and
- $\mathcal{B}$  is the path that remains after removing the path from the first return of the path from the right and the Motzkin path from  $y$  to  $x$  (including  $x$ ) from the original path.

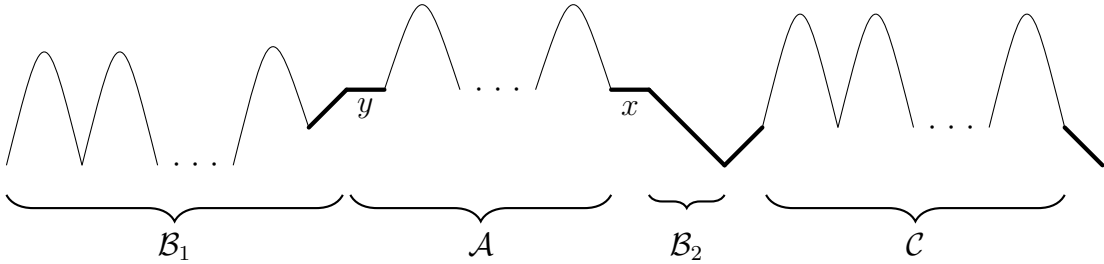


FIGURE 1. Canonical decomposition

This process is performed recursively and terminates at an empty path. Note that each application of  $\Phi$  adds one node and removes one of each type of step. This proves inductively that an S-Motzkin path of length  $3n$  maps to a ternary tree with  $n$  nodes.

2.2. **Ternary trees to S-Motzkin paths.** The inverse mapping is performed recursively bottom-up as follows. Each node of a ternary tree has three subtrees. Call the paths associated with the left, middle, and right subtrees  $\mathcal{A}$ ,  $\mathcal{B}$ , and  $\mathcal{C}$  respectively.

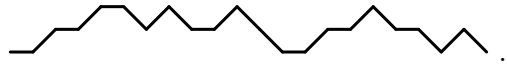
Starting at the end nodes, replace each node with

$$\mathcal{B}_1 \mathcal{A} \_ \mathcal{B}_2 / \mathcal{C} \setminus,$$

where  $\mathcal{B}_1$  is the subpath of  $\mathcal{B}$  that starts at  $(0,0)$  and extends to and includes the first occurring  $/$  from the right. The path  $\mathcal{B}_2$  is what remains of  $\mathcal{B}$  after removing  $\mathcal{B}_1$ .

This process is continued recursively on each set of end nodes and terminates at the root to produce an S-Motzkin path. Note that for each node three steps are added, and thus a ternary tree with  $n$  nodes produces an S-Motzkin path of length  $3n$ .

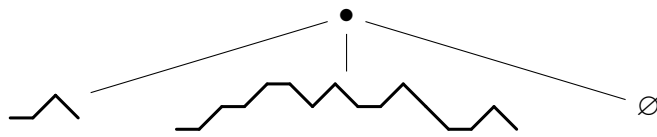
2.3. **Example.** As an example, we map the following S-Motzkin path into a ternary tree. Since the steps are reversible, the inverse mapping can be seen by reading the example in reverse. Let  $\mathcal{M}$  be



The canonical decomposition of  $\mathcal{M}$  is

$$\Phi(\mathcal{M}) = (\text{short path}, \text{long path}, \emptyset).$$

Hence



Continuing recursively:

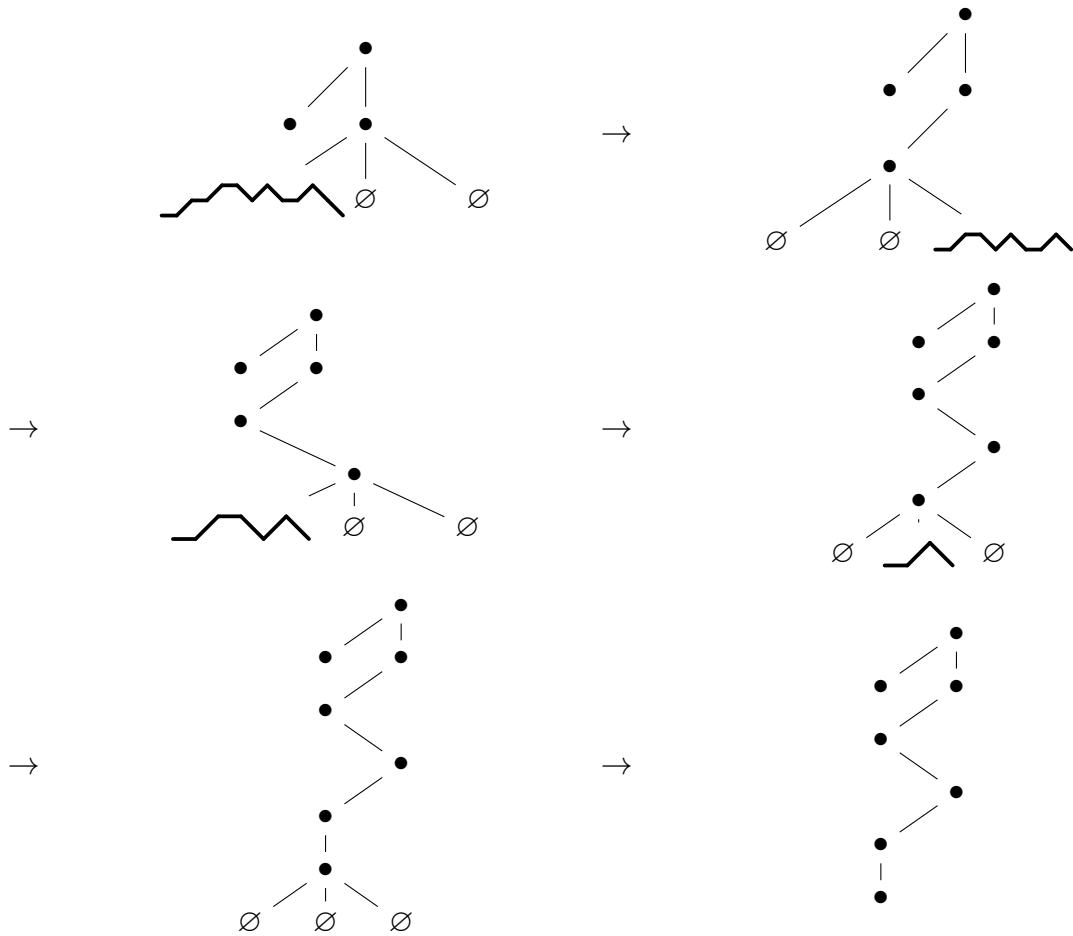

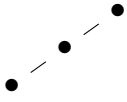


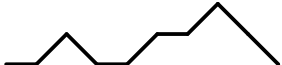





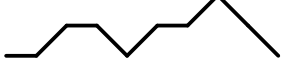

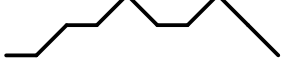

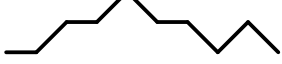

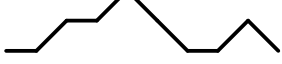



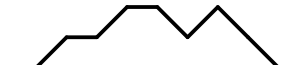

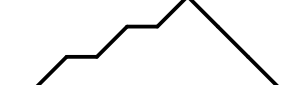



TABLE 1. Bijection for  $n = 3$ 

S-Motzkin path	Ternary tree
	
	
	
	
	
	
	
	
	
	
	
	

## 3. GENERALIZATION

This bijection can be generalized to be between  $t$ -ary trees and the subclass of Motzkin paths with  $(t - 2)n$   $\_$  steps and  $n$  of each of the other steps such that

- The initial  $t - 2$  steps must be of the form  $\_$ ,
- between every two  $\_$  there are exactly  $t - 2$  steps of the form  $\_$ ,
- the  $k$ -th occurring  $\backslash$  must occur after at least  $k$  occurrences of  $t - 2$  steps of the form  $\_$  and one step of the form  $\_$ .

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