

# Signed $\beta$ -expansions of minimal weight

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## Expansions in base 2

Every integer  $N \geq 0$  has an expansion in base 2

$$N = \sum_{j=0}^K \epsilon_j 2^j = \epsilon_K \cdots \epsilon_1 \epsilon_0.$$

with  $\epsilon_j \in \{0, 1\}$ , which is unique up to leading zeros.

If we allow negative digits, then the number of non-zero digits can often be reduced:

$$7 = 4 + 2 + 1 = 111. = 100\bar{1}. = 8 - 1 \quad (\bar{1} = -1)$$

**Problem:** find an expansion of  $N$  of minimal weight  $\sum_{j=0}^K |\epsilon_j|$

(cf. Hamming weight: number of non-zero digits  $\epsilon_j$ , equal to this weight if  $\epsilon_j \in \{-1, 0, 1\}$ )

# Expansions of minimal weight in base 2

Booth (1951), Reitwiesner (1960), . . . : **Non-Adjacent Form (NAF)**

Every integer  $N$  has a unique expansion  $N = \sum_{j=0}^K \epsilon_j 2^j$  with  $\epsilon_j \in \{-1, 0, 1\}$  such that  $\epsilon_{j-1} = \epsilon_{j+1} = 0$  if  $\epsilon_j \neq 0$ . The weight of this expansion is minimal among all expansions of  $N$  in base 2.

Dajani, Kraaikamp, Liardet (2006): ergodic properties of the dynamical system associated with the NAF,

$$T : [-2/3, 2/3) \rightarrow [-2/3, 2/3), \quad T(x) = 2x - \lfloor (3x + 1)/2 \rfloor$$

Heuberger (2004):  $\epsilon_K \cdots \epsilon_1 \epsilon_0 \in \{\bar{1}, 0, 1\}^*$  is a signed 2-expansion of minimal weight if and only if contains none of the factors

$$11(01)^*1, \quad 1(0\bar{1})^*\bar{1}, \quad \bar{1}\bar{1}(0\bar{1})^*\bar{1}, \quad \bar{1}(01)^*1.$$

( $A^*$  is the free monoid over the set  $A$ ,

$$a^* = \{a\}^* = \{\text{empty word}, a, aa, aaa, aaaa, \dots\})$$

joint digit expansions: Solinas; Grabner, Heuberger, Prodinger

$\beta$ -expansions,  $\beta = \frac{1+\sqrt{5}}{2}$

Greedy  $\beta$ -expansions: Every  $x \in \mathbb{R}^+$  has a unique expansion

$$x = \sum_{j \in \mathbb{Z}} \epsilon_j \beta^{-j} = \cdots \epsilon_{-1} \epsilon_0 . \epsilon_1 \epsilon_2 \cdots$$

with  $\epsilon_j \in \{0, 1\}$ ,  $\epsilon_{j-1} = \epsilon_{j+1} = 0$  if  $\epsilon_j = 1$ , which does not terminate with  $(10)^\omega = 101010 \cdots$ .

$$\beta^2 = \beta + 1, \quad 100. = 011., \quad 1. = .11$$

Greedy  $\beta$ -expansions are not minimal in weight for  $\epsilon_j \in \{-1, 0, 1\}$ :

$$0101001. = 10\bar{1}1001. = 1000\bar{1}01. = 10000\bar{1}0.$$

## $\beta$ -expansions of minimal weight

$x = x_1 \cdots x_n \in A_\beta^*$  is  $\beta$ -heavy if it is not minimal in weight, i.e., if there exists  $y = y_\ell \cdots y_r \in A_\beta^*$  with

$$\bullet x_1 \cdots x_n = y_\ell \cdots y_0 \bullet y_1 \cdots y_r \quad \text{and} \quad \sum_{j=\ell}^r |y_j| < \sum_{j=1}^n |x_j|.$$

If  $x_1 \cdots x_{n-1}$  and  $x_2 \cdots x_n$  are not  $\beta$ -heavy,  $x$  is strictly  $\beta$ -heavy.

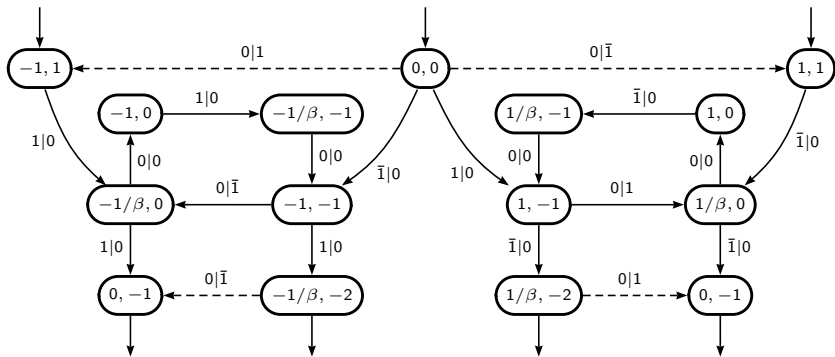
### Theorem

If  $\beta = \frac{1+\sqrt{5}}{2}$ , then the set of strictly  $\beta$ -heavy words is

$$\begin{aligned} &1(0100)^*1 \cup 1(0100)^*0101 \cup 1(00\bar{1}0)^*\bar{1} \cup 1(00\bar{1}0)^*0\bar{1} \\ &\cup \bar{1}(0\bar{1}00)^*\bar{1} \cup \bar{1}(0\bar{1}00)^*0\bar{1}0\bar{1} \cup \bar{1}(0010)^*1 \cup \bar{1}(0010)^*01. \end{aligned}$$

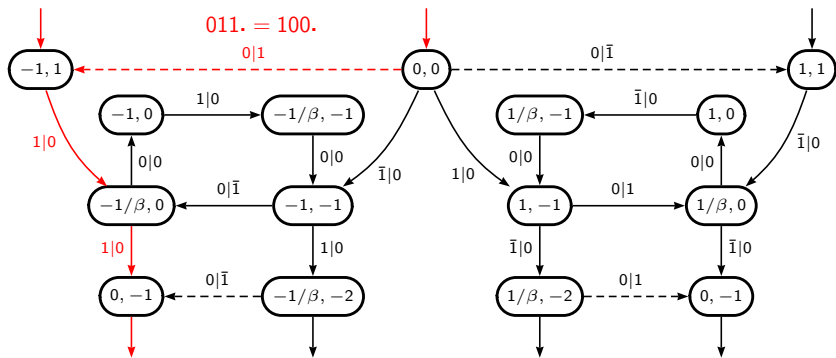
If  $\cdots \epsilon_{-1} \epsilon_0 \epsilon_1 \cdots$  does not contain any of these factors, then  $\cdots \epsilon_{-1} \epsilon_0 \bullet \epsilon_1 \cdots$  is a signed  $\beta$ -expansion of minimal weight.

The strictly  $\beta$ -heavy words are the inputs of the following transducer. The outputs are corresponding lighter words (if the path is completed by dashed arrows such that it runs from  $(0, 0)$  to  $(0, -1)$ ).



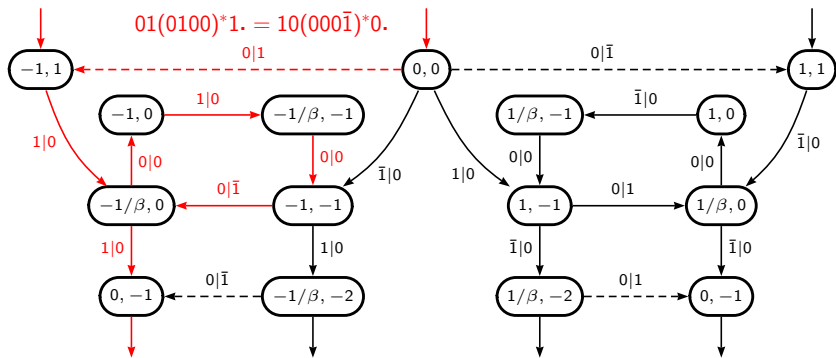
$$(s, \delta) \xrightarrow{a|b} (s', \delta') : s' = \beta s + a - b, \delta' = \delta + |b| - |a|$$

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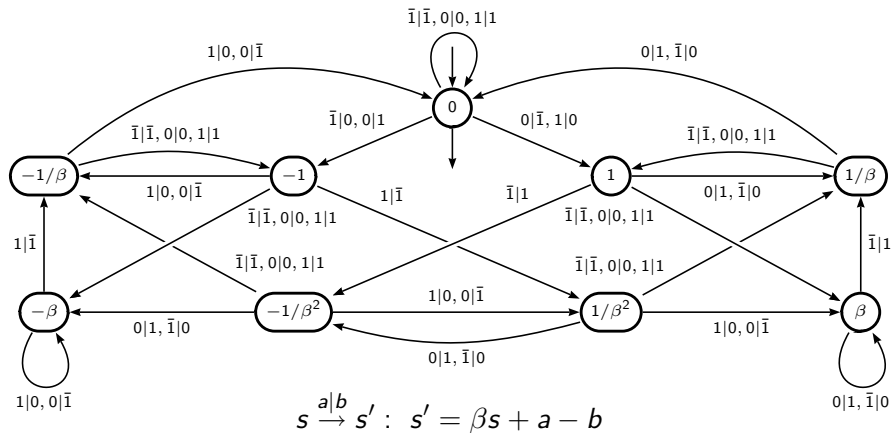
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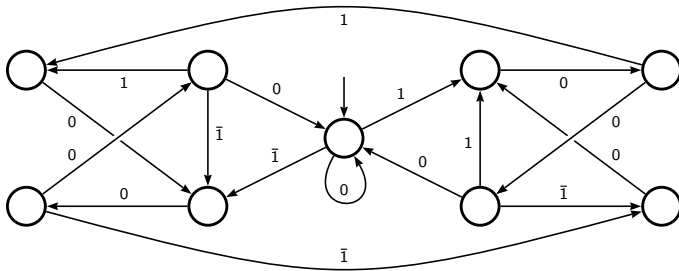
The sequences  $x_1|y_1, \dots, x_n|y_n$  with  $x_1 \cdots x_n, y_1 \cdots y_n \in \{\bar{1}, 0, 1\}^*$  (not containing a factor  $11$  or  $\bar{1}\bar{1}$ ) such that  $\cdot x_1 \cdots x_n = \cdot y_1 \cdots y_n$  are recognized by the redundancy automaton (transducer)



If  $s_0 = 0$ ,  $s_{j-1} \xrightarrow{x_j|y_j} s_j$ ,  $1 \leq j \leq n$ , then  $s_j = x_1 \cdots x_j \cdot - y_1 \cdots y_j \cdot$ , and  $\cdot x_1 \cdots x_n = \cdot y_1 \cdots y_n$  if and only if  $s_n = 0$ .

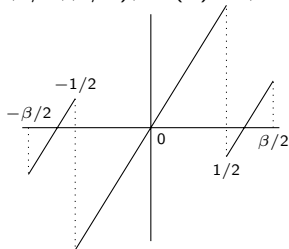
## Theorem

For  $\beta = \frac{1+\sqrt{5}}{2}$ , the signed  $\beta$ -expansions of minimal weight are given by the following automaton, where all states are terminal.



# Transformation providing signed $\beta$ -expansions of minimal weight, $\beta = \frac{1+\sqrt{5}}{2}$

$$T : [-\beta/2, \beta/2) \rightarrow [-\beta/2, \beta/2), \quad T(x) = \beta x - \lfloor x + 1/2 \rfloor$$



## Proposition

*If  $x \in [-\beta/2, \beta/2)$  and  $x_j = \lfloor T^{j-1}(x) + 1/2 \rfloor$ , then  $x = .x_1 x_2 \dots$  is a signed  $\beta$ -expansion of minimal weight avoiding the factors 11, 101,  $1\bar{1}$ ,  $10\bar{1}$ ,  $100\bar{1}$  and their opposites.*

## Proof of the proposition.

Recall that  $T(x) = \beta x - \lfloor x + 1/2 \rfloor$  and  $x_j = \lfloor T^{j-1}(x) + 1/2 \rfloor$ .  
If  $x_j = 1$ , then  $T^{j-1}(x) \in [1/2, \beta/2) = [.(010)^\omega, .(100)^\omega)$ ,

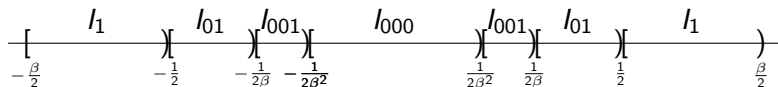
$$\begin{aligned} T^j(x) &\in [\beta/2 - 1, \beta^2/2 - 1) = [-1/(2\beta^2), 1/(2\beta)), & x_{j+1} &= 0, \\ T^{j+1}(x) &\in [-1/(2\beta), 1/2), & x_{j+2} &= 0, \\ T^{j+2}(x) &\in [-1/2, \beta/2), & x_{j+3} &\in \{0, 1\}. \end{aligned}$$

Hence  $11$ ,  $101$ ,  $1\bar{1}$ ,  $10\bar{1}$  and  $100\bar{1}$  are avoided, thus the strictly  $\beta$ -heavy words  $1(0100)^*1$ ,  $1(0100)^*0101$ ,  $1(00\bar{1}0)^*\bar{1}$ ,  $1(00\bar{1}0)^*0\bar{1}$  are avoided. The same is true for the opposite words.  $\square$

**Remark.** Heuberger (2004) excluded (for the Fibonacci numeration system) the factor  $1001$  instead of  $100\bar{1}$ . This can be achieved by  $T(x) = \beta x - \lfloor \frac{\beta^2+1}{2\beta}x + \frac{1}{2} \rfloor$  on  $[\frac{-\beta^2}{\beta^2+1}, \frac{\beta^2}{\beta^2+1})$ ,  $\frac{\beta^2}{\beta^2+1} = .(1000)^\omega$ .

# Markov chain of digits

Let  $T(x) = \beta x - \lfloor x + 1/2 \rfloor$ , and  $I_{000}, I_{001}, I_{01}, I_1$  as follows



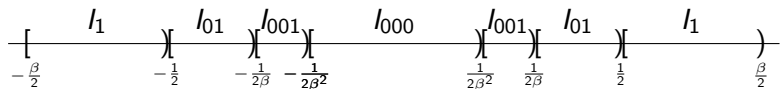
The sequence of random variables  $(X_k)_{k \geq 0}$  defined by

$$\begin{aligned} \Pr[X_0 = j_0, \dots, X_k = j_k] \\ &= \lambda(\{x \in [-\beta/2, \beta/2) : x \in I_{j_0}, T(x) \in I_{j_1}, \dots, T^k(x) \in I_{j_k}\})/\beta \\ &= \lambda(I_{j_0} \cap T^{-1}(I_{j_1}) \cap \dots \cap T^{-k}(I_{j_k}))/\beta \end{aligned}$$

(where  $\lambda$  denotes the Lebesgue measure) is a Markov chain since

$$T(I_{000}) = I_{000} \cup I_{001} = T(I_1), \quad T(I_{001}) = I_{01}, \quad T(I_{01}) = I_1$$

and  $T(x)$  is linear on each  $I_j$ .



The matrix of transition probabilities is

$$(\Pr[X_k = j \mid X_{k-1} = i])_{i,j \in \{000,001,01,1\}} = \begin{pmatrix} 1/\beta & 1/\beta^2 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 2/\beta^2 & 1/\beta^3 & 0 & 0 \end{pmatrix}$$

the stationary distribution vector is  $(2/5, 1/5, 1/5, 1/5)$ . Therefore

$$\Pr[X_k = 1] = \lambda(\{x \in [-\beta/2, \beta/2] : T^k(x) \in I_1\}) \rightarrow 1/5,$$

i.e., the expected number of non-zero digits in a signed  $\beta$ -expansion of minimal weight of length  $n$  is  $n/5 + \mathcal{O}(1)$ .  
(cf. greedy  $\beta$ -expansions  $n/(\beta^2 + 1)$ , base 2 minimal expansions  $n/3$ )

# Fibonacci numeration system

Let  $F_0 = 1$ ,  $F_1 = 2$ ,  $F_j = F_{j-1} + F_{j-2}$ . Then every integer  $N \geq 0$  has a unique  $F$ -expansion

$$N = \sum_{j=1}^n \epsilon_j F_{n-j} = \langle \epsilon_1 \cdots \epsilon_n \rangle_F$$

with  $\epsilon_j \in \{0, 1\}$  and  $\epsilon_{j-1} = \epsilon_{j+1} = 0$  if  $\epsilon_j = 1$ .

$x_1 \cdots x_n \in \{\bar{1}, 0, 1\}^*$  is  **$F$ -heavy** if there exists  $y_\ell \cdots y_n \in \{\bar{1}, 0, 1\}^*$  such that  $\langle x_1 \cdots x_n \rangle_F = \langle y_\ell \cdots y_n \rangle_F$  and  $\sum_{j=\ell}^n |y_j| < \sum_{j=1}^n |x_j|$ .  
 $\langle \cdots 1\bar{1}0 \cdots \rangle_F = \langle \cdots 001 \cdots \rangle_F$ , but  $\langle \cdots 1\bar{1} \rangle_F = \langle \cdots 01 \rangle_F$ .

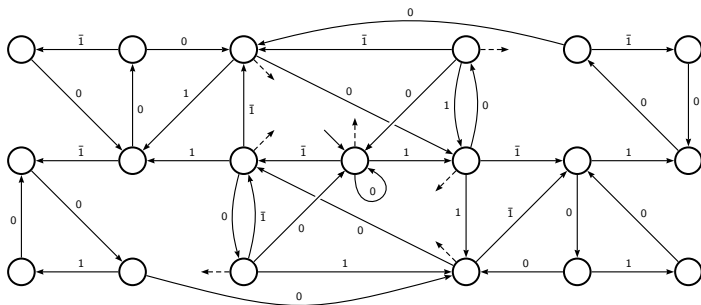
## Theorem

*The  $F$ -heavy words are exactly the  $\beta$ -heavy words for  $\beta = \frac{1+\sqrt{5}}{2}$ , i.e. a word is a signed  $F$ -expansion of minimal weight if and only if it is a signed  $\beta$ -expansion of minimal weight.*

## Tribonacci numeration system

If  $\beta > 1$  is the Tribonacci number,  $\beta^3 = \beta^2 + \beta + 1$ , then the strictly  $\beta$ -heavy words are the inputs of the automaton on the next slide. The signed  $\beta$ -expansions of minimal weight are given by the following automaton where all states are terminal.

If  $T_0 = 1$ ,  $T_1 = 2$ ,  $T_2 = 4$ ,  $T_j = T_{j-1} + T_{j-2} + T_{j-3}$ , then the signed  $T$ -expansions of minimal weight are given by the automaton where only the states with dashed outgoing arrows are terminal.







## Particular signed $\beta$ -expansions of minimal weight, $\beta^3 = \beta^2 + \beta + 1$

$$T : \left[ \frac{-\beta^2}{\beta^2+1}, \frac{\beta^2}{\beta^2+1} \right) \rightarrow \left[ \frac{-\beta^2}{\beta^2+1}, \frac{\beta^2}{\beta^2+1} \right), \quad T(x) = \beta x - \left\lfloor \frac{\beta^2+1}{2\beta} x + \frac{1}{2} \right\rfloor$$

### Proposition

If  $x \in \left[ \frac{-\beta^2}{\beta^2+1}, \frac{\beta^2}{\beta^2+1} \right)$  and  $x_j = \left\lfloor \frac{\beta^2+1}{2\beta} T^{j-1}(x) + 1/2 \right\rfloor$ , then  $x = \cdot x_1 x_2 \cdots$  is a signed  $\beta$ -expansion of minimal weight avoiding the factors  $11$ ,  $1\bar{1}$ ,  $10\bar{1}$  and their opposites.

The transition probabilities of the corresponding Markov chain are

$$\begin{pmatrix} 1/\beta & 1 - 1/\beta & 0 \\ 0 & 0 & 1 \\ 1 - 1/\beta^2 & 1/\beta^2 & 0 \end{pmatrix}$$

the stationary distribution vector is  $\left( \frac{\beta^3+\beta^2}{\beta^5+1}, \frac{\beta^3}{\beta^5+1}, \frac{\beta^3}{\beta^5+1} \right)$ , hence the probability of a non-zero digit is asymptotically  $\frac{\beta^3}{\beta^5+1} \approx 0.28219$ .

## General case, condition (D)

(D):  $\beta > 1$  and  $P(\beta) = 0$  for some polynomial  
 $P(X) = X^d - b_1 X^{d-1} - \dots - b_d \in \mathbb{Z}[X]$  with  $b_1 > \sum_{j=2}^d |b_j|$

If  $\beta$  satisfies (D), then  $\beta$  is a Pisot number.

### Proposition (Akiyama, Rao, St. (2004))

*Let  $\beta$  satisfy (D), and  $x_1 \cdots x_n \in \mathbb{Z}^*$  such that  $|\cdot x_1 \cdots x_n| < 1$ .*

*Then there exists a word  $y_0 \cdots y_m \in \{-\lfloor \beta \rfloor, \dots, \lfloor \beta \rfloor\}^*$  such that  $y_0 \cdot y_1 \cdots y_m = \cdot x_1 \cdots x_n$  and  $\sum_{j=0}^m |y_j| \leq \sum_{j=1}^n |x_j|$ .*

(More precisely, the positive and negative parts of  $y_0 \cdots y_m$  can be chosen to be  $\beta$ -admissible, hence  $\beta$  satisfies (W).)

### Theorem

*If  $\beta$  satisfies (D), then the set of signed  $\beta$ -expansions of minimal weight is recognized by a finite automaton, which is computable.*